

# Semantic Program Optimization

## Avoiding (some) Data Dependencies

### Introduction

- **Tiling** is an effective transformation for parallelism, locality improvement and granularity.
- However, its legality relies on not violating **program dependencies**.

- We propose to use **semantic properties** (associativity, commutativity) to do tiling.
- The semantically-equivalent tiled program has **different dependencies**.
- The derivation of this transformation relies on **program equivalence checking**.

### Using Semantic Properties for Tiling

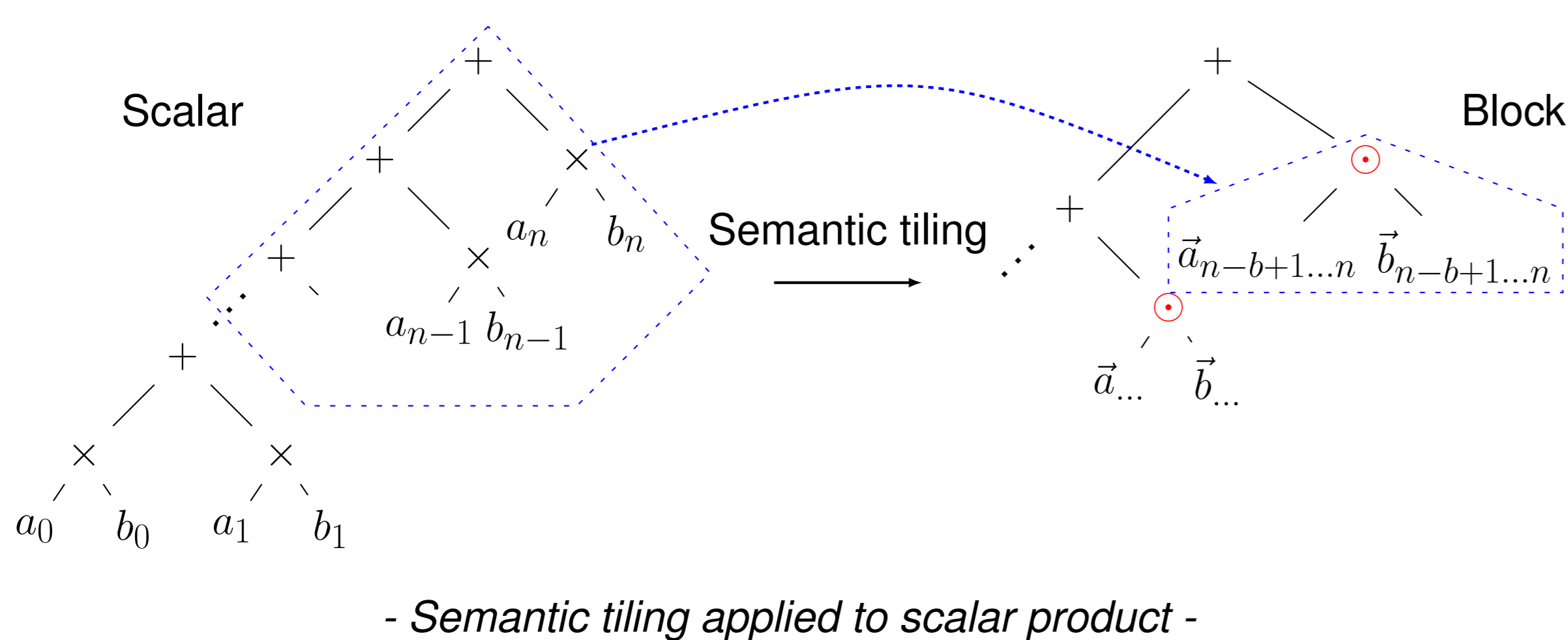
**Semantic tiling:** (for this poster) Transformation of a program operating on **scalars** into an equivalent program operating on **blocks** with the same structure.

- **"Semantic":** Operators are not the same (semantic rules are used)
- **"Tiling":** Operands on blocks instead of scalars (raise the **level of abstraction**)  
⇒ Locality benefits (like in classical tiling).

**Properties of semantic tiling:**

- **Reorganize operations** and dependencies of a program
- **Recursive call** on a smaller instance → Divide-and-conquer scheme.
- Apply on **linear algebra** and **graph theoretic** algorithms (need a notion of block)
- Applicability conditions are **different** than for classical tiling.  
(semantic tilable but non-classically tilable: Algebraic Path Problem)

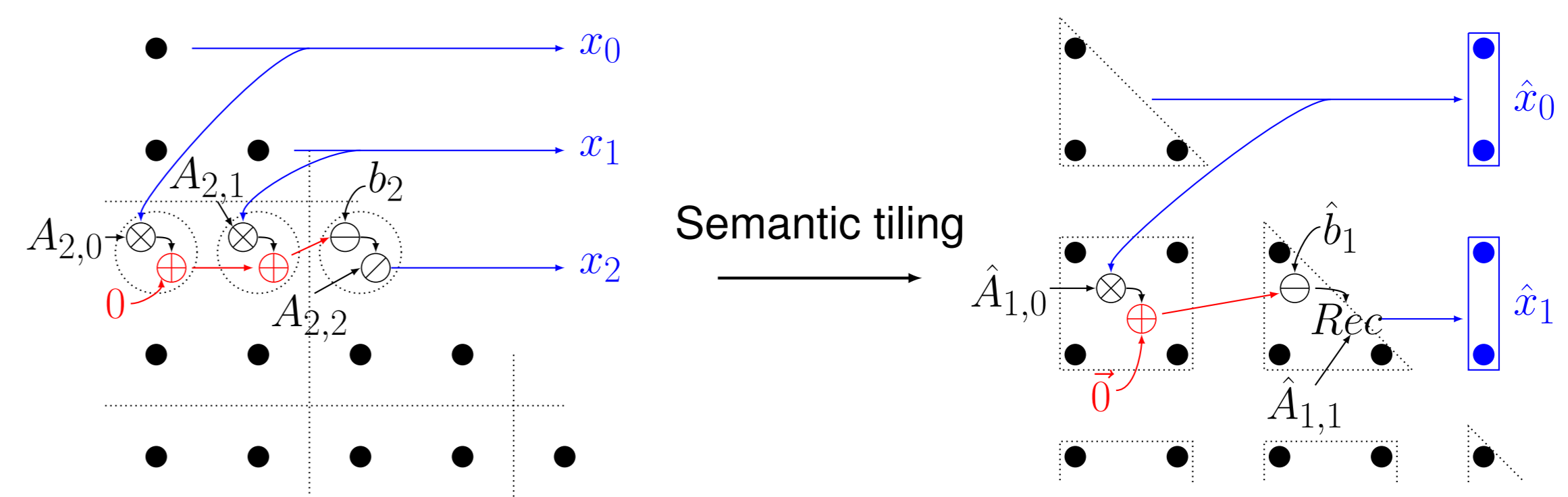
**Semantically tilable examples:** **Matrix multiply**, **Forward substitution**, **LU**, **Cholesky**, sub-problems of **APP** (**shortest path** in a graph), ...



### Toward Automatic Derivation

Two approaches are being explored:

1. Using sophisticated version of "pattern matching" to **hypothesize the structure of the blocked program**. Then, **proving its validity** through program equivalence checking.
2. **Transforming** the scalar program into its semantic tiled version, by using **rewriting rules** based on algebraic properties. More precisely:
  - a. **Data tiling** of the inputs, outputs and temporaries of the program.
  - b. **Extracting the program slice** that touches the same data block.
  - c. **Check the equivalence** with a corresponding matrix-level operator.
  - d. **Substituting** them by the matrix-level operators → Semantic tiled program.



- Semantic tiling applied to forward substitution -

### Checking Program Equivalence

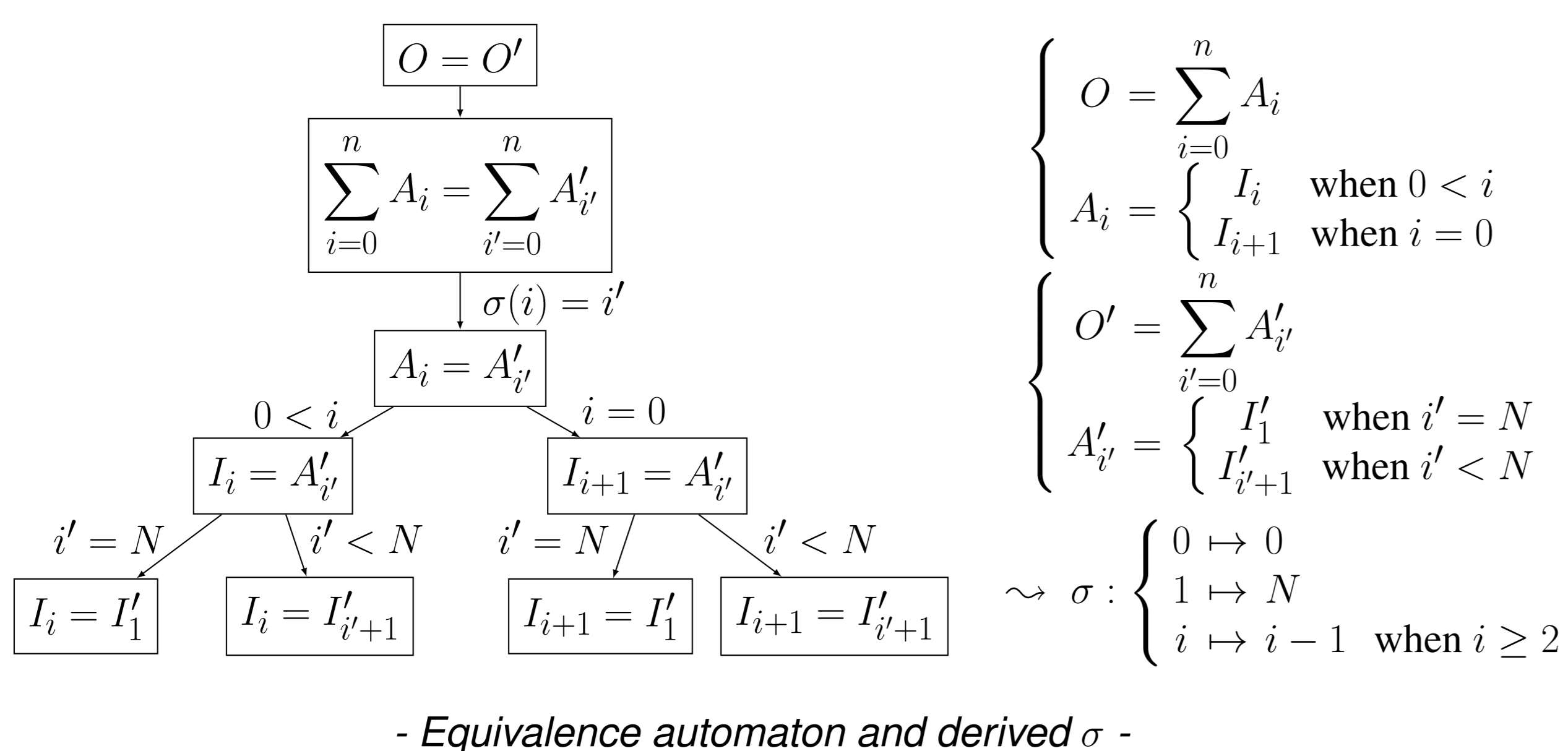
We build on the 2-step **equivalence algorithm** for SAREs proposed on [1]:

- Compile the program unification problem as an **integer interpreted automaton**.
- Then, check **reachability** by **symbolic execution** of this automaton.

Do not manage any semantic rules (such as **associativity** and **commutativity**).

⇒ We proposed [2] an extension of the equivalence algorithm to manage **reductions**, i.e. successive applications of an associative and commutative operator over a family of expressions (example:  $\sum_{i=0}^{n-1} A[i]$ ).

**Main idea:** find a **bijection**  $\sigma$  between the reduction bodies.



### Conclusion

**Semantic tiling** is a **semantic program transformation** that raises the level of abstraction of a program.

⇒ Allows us to **"break"** some dependencies during tiling, while retaining program semantics.

⇒ **Raises the granularity** of the considered operations → Suitable for HLS.

### Future Work

**Semantic tiling automatic derivation:**

- Mathematical formalization in progress. Encountered issues:
  - How to systematically group instructions together?
  - How to prune the higher-order operator?
  - How do we recognize an instruction block corresponding to a recurrence call?
- How to characterize the programs that can be semantically tiled?
- Generalize beyond just blocking.

**Program equivalence with reductions:**

- Implementation in progress.
- Improve the applicability of the equivalence algorithm.
- Can be reused for other purpose (transformation/validation)

### References

- [1] *On the Equivalence of Two System of Affine Recurrence Equations* (by D.Barthou, P.Feautrier and X.Redon)  
[2] *On Program Equivalence with Reductions* (by G.Iooss, C.Alias, S.Rajopadhye, technical report in preparation)

### People Involved

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